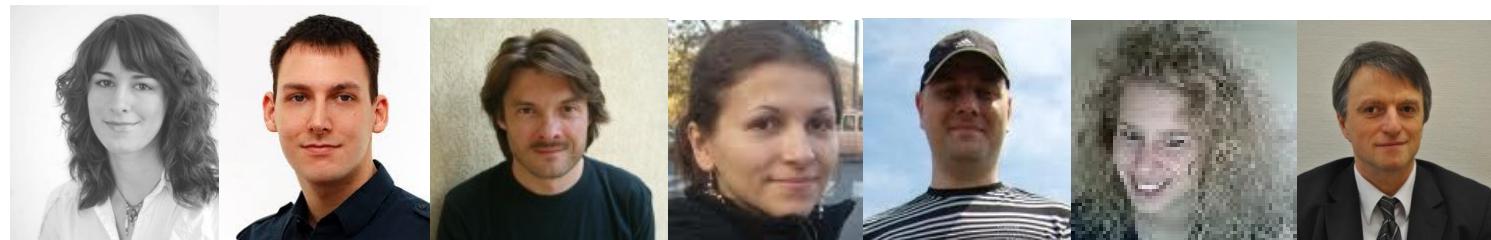




# Parallelization by Refactoring

Dept. Programming Languages and Compilers  
Eötvös Loránd University, Hungary



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# LET'S PARTY!

- Have fun...

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# LET'S PaRTE!

- Have fun...

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# LET'S PaRTE!

- Have functional programming!

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# LET'S PaRTE!

- Have functional programming!

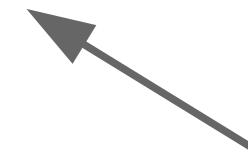
map-like function                  speedup prediction  
**pattern candidate discovery**  
static analysis                  task farm  
refactoring    pipeline      parallel patterns  
algorithmic skeletons      divide and conquer  
**ParaPhrase Refactoring Tool for Erlang**  
RefactorErl                  Wrangler

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# Motivation

- Highly heterogeneous mega-core computers
- Performance and energy
- Think parallel
  - High-level programming constructs
  - Deadlocks etc. eliminated by design
  - Communication packaged/abstracted
  - Performance information is part of design
- Restructure existing code



Kevin  
Hammond



# How shall I parallelize?

- Refactor
  - Use a tool!
  - Guided, semi-automatic transformations
- Experiment
  - Measure, validate
- Repeat
- Applicable for legacy code as well

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# Where shall I parallelize?

- Independent computations
- Good potential for speedup
  - Complex computation?
  - Low sequential overhead?
- Find candidates automatically
  - Use a tool!
  - Static program analysis

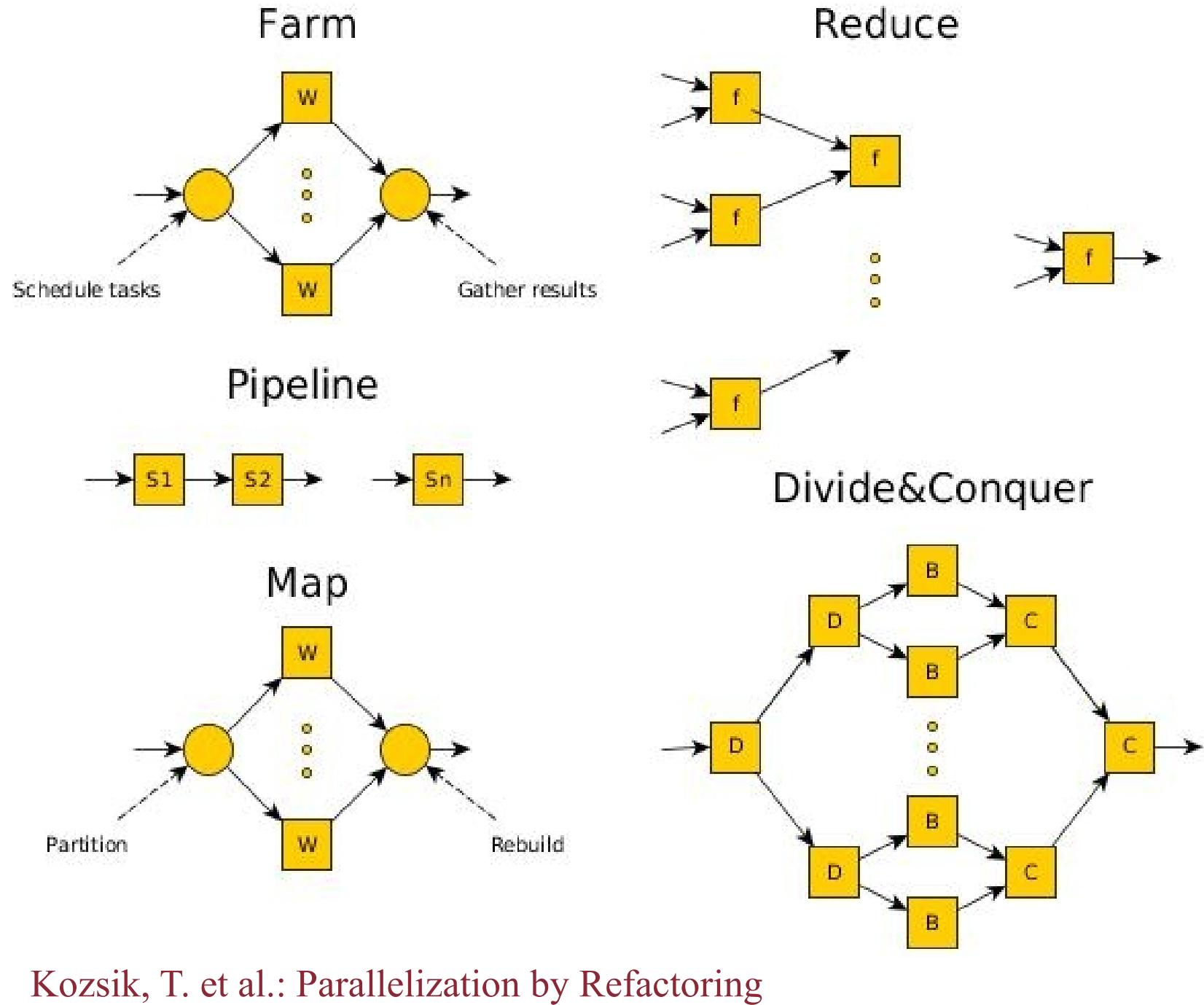


# Pattern-based parallelism

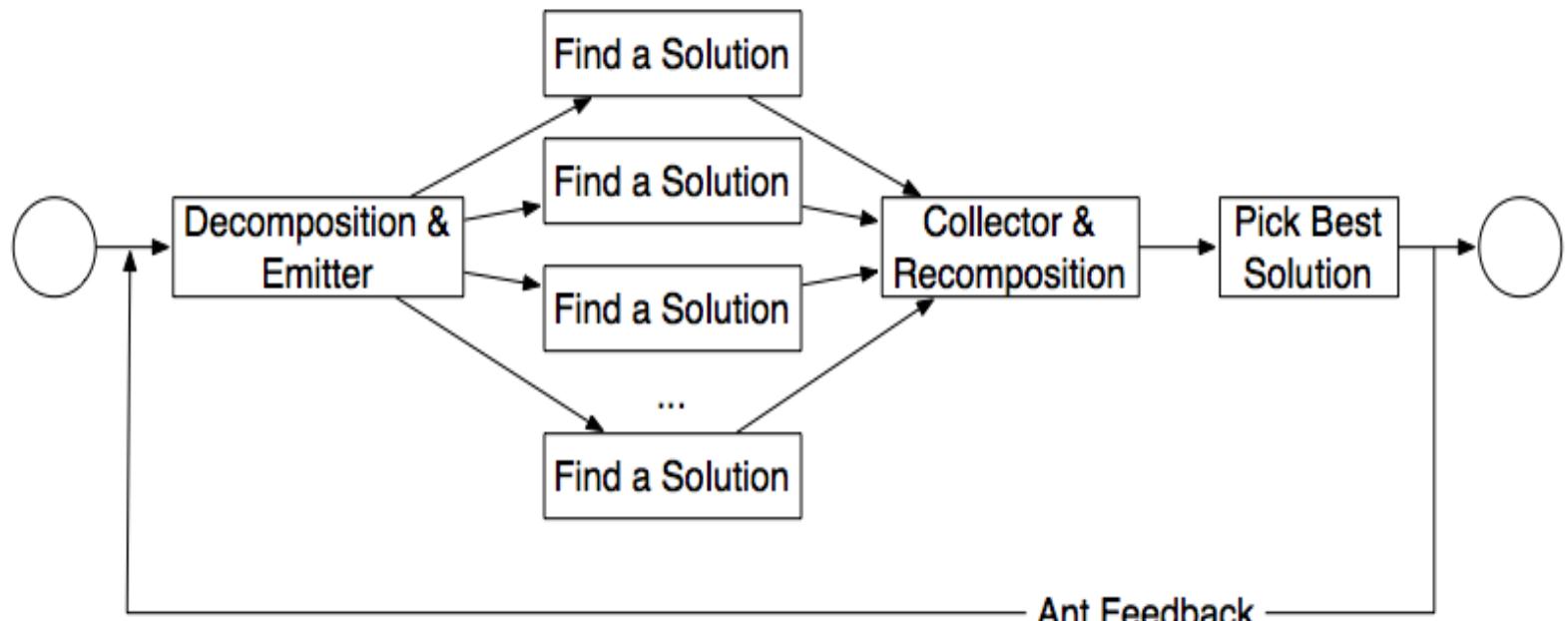
High-level approach to parallel programming

- Rely on a library of algorithmic skeletons
- Easier to develop code
- Easier to modify / maintain
- Better utilization of resources
  - Static resource management
  - Dynamic resource management

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# Pool pattern



## Ant Colony Optimization

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## Multithreaded programming



picture from Kevin Hammond

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# Parallel Patterns for Adaptive Heterogeneous Multicore Systems

- Programmability of heterogeneous parallel architectures
- Structured design and implementation of parallelism
- High-level parallel patterns
- Dynamic (re)mapping on heterogeneous hardware

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# PARAPHRASE



University  
of  
St Andrews



ROBERT GORDON  
UNIVERSITY • ABERDEEN



Queen's University  
Belfast

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DEGLI STUDI  
DI TORINO  
  
ALMA UNIVERSITAS  
TAURINENSIS



UNIVERSITÀ DI PISA

Universität Stuttgart



National  
College of  
Ireland

Cloud Competency Center



Erlang  
SOLUTIONS



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- functional language
- strict, impure, dynamically typed
- concurrency, actor model
- distribution
- fault tolerance
- Open Telecom Platform

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# Companies using Erlang



ERICSSON

amazon.com<sup>®</sup>

Goldman  
Sachs



SMARKETS

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THE HUFFINGTON POST  
THE INTERNET NEWSPAPER. NEWS BLOGS VIDEO COMMUNITY

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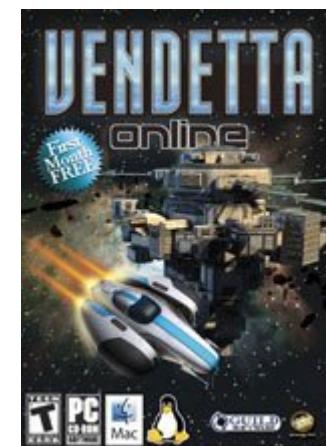
# Products using Erlang



## GitHub



## CALL OF DUTY



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# Let's PaRTE!

Assist parallelization of Erlang code with the

## ParaPhrase Refactoring Tool for Erlang

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# So what is Erlang like?

- Functional (strict, impure, dynamically typed)
- Garbage collected
- Designed for concurrency and distribution
  - Lightweight processes
  - Actor model (message passing)
- Processing binary data
- Fault tolerance (failure recovery)
- Compiles to **beam**, runs in VM
- Hot code swap

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# Erlang: a functional language

- Variables bound only once
- Recursion instead of loops
  - Tail recursive
- Higher-order functions instead of recursion
- Lambdas, pattern matching
- Properties
  - Strict evaluation (call-by-value)
  - Impurity (e.g. communication)
  - No partial applications



# Terms

- Literals, e.g. 42 or 42.0
- Atoms, e.g. leaf, blue, ok, error, true
- Compound data

Funs	<code>fun(X) -&gt; X+1 end</code>
Lists	<code>[0,1,1,2,3,5,8,13,21]</code>
Tuples	<code>{may, 10, 2014}</code>
Records	<code>#date{month=june,day=7,year=2015}</code>
Binaries	<code>&lt;&lt;0,1,1,2,3,5,8,13,21&gt;&gt;</code>

- Pids, ports, refs



# A complex term

```
[  
 {  
   farm,  
   [  
     {  
       pipe,  
       [  
         {seq, fun scan/1},  
         {seq, fun parse/1}  
       ]  
     }  
   ],  
   12  
 }]  
 ]
```

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# Expressions

- Terms (literals, atoms, compound data)
- Variables, e.g. `X`, `Long_Variable_Name`
- Function/operator calls, e.g. `fib(N-1)+fib(N-2)`
- Data structures, e.g. `{june,Day,fib(18)-569}`
- Control structures
  - Branching (**case** and **if**)
  - Sending and receiving a message
  - Error handling



# Defining a named function

increment(N) -> N+1.



# Pattern matching

```
fib(N) ->
```

```
  case N of
```

```
    0 -> 0;
```

```
    1 -> 1;
```

```
    _ -> fib(N-1) + fib(N-2)
```

```
  end.
```



# Multiple function clauses

```
fib(0) -> 0;
```

```
fib(1) -> 1;
```

```
fib(N) -> fib(N-1) + fib(N-2).
```



# Guards

```
fib(N) when N < 2 -> N;  
fib(N) -> fib(N-1) + fib(N-2).
```



# Using an if-expression

```
fib(N) ->  
  if  
    N < 2 -> N;  
    true -> fib(N-1) + fib(N-2)  
  end.
```



# Recursion

```
factorial(1) -> 1;  
factorial(N) -> N * factorial(N-1).
```

---



# Tail recursion

factorial(1) -> 1;  
factorial(N) -> N \* factorial(N-1).



factorial(N) -> factorial\_acc(N, 1).  
factorial\_acc(1, Acc) ->  
Acc;  
factorial\_acc(N, Acc) ->  
**factorial\_acc(N-1, Acc\*N)**.

# Overloading on arity

```
prime(1) -> false;  
prime(N) when N > 1 -> prime(N,2).
```

% no proper divisors of  $N$   
% between  $M$  and  $\text{sqrt}(N)$

```
prime(N,M) ->  
    M*M>N orelse ( N rem M =/= 0 andalso  
                      prime(N,M+1)  
                    ).
```

- Overloaded functions `prime/1` and `prime/2`
- `prime/2` is helper function



# Lists

[0,1,1,2,3,5,8]

[0 | [1,1,2,3,5,8]]

[0 | [1 | [1 | [2 | [3 | [5 | [8 | []]]]]]]]]

[0,1,1,2 | [3 | [5 | [8 | []]]]]]

[0,1,1,2 | [3,5,8]]

- [Head | Tail] notation
- List comprehension

[N | | N <- lists:seq(1,100), prime(N)]



# Scalar-vector multiplication

```
mul(Scalar, []) -> [];  
mul(Scalar, [Head|Tail]) ->  
[ Scalar*Head | mul(Scalar, Tail) ].
```



# Scalar-vector multiplication

```
mul(Scalar,[]) -> [];  
mul(Scalar,[Head|Tail]) ->  
  [ Scalar*Head | mul(Scalar,Tail) ].
```

---

```
mul(Scalar,List) ->  
  [ Scalar*Item || Item <- List ].
```

# Scalar-vector multiplication

```
mul(Scalar,[]) -> [];
```

```
mul(Scalar,[Head|Tail]) ->  
[ Scalar*Head | mul(Scalar,Tail) ].
```

---

```
mul(Scalar,List) ->  
[ Scalar*Item || Item <- List ].
```

---

```
mul(Scalar,List) ->  
map( fun(Item) -> Scalar*Item end, List ).
```



# Higher-order functions

```
map(Fun, []) -> [];  
map(Fun, [Head|Tail]) ->  
  [ Fun(Head) | map(Fun, Tail) ].
```

```
filter(Pred,List) ->  
  [ Item || Item <- List, Pred(Item) ]
```





# fun-expressions

```
map( fun(Item) -> Scalar*Item end, List )
```

```
filter( fun prime/1, List )
```



# Variable binding

- Formal parameters:

`fib(N) -> ...`

`mul(Scalar, [Head|Tail]) -> ...`

- Generator in list comprehension:

`[ ... | Item <- List ]`

- Pattern matching expression

Syntax: Pattern = Expression

Primes = primes(List)

[ Head | Tail ] = primes(List)



# Sequence of expressions

```
area( {square, Side} ) ->  
    Side * Side;
```

```
area( {circle, Radius} ) ->  
    % almost :-)  
    3.14 * Radius * Radius;
```

```
area( {triangle, A, B, C} ) ->  
    S = (A + B + C)/2,  
    math:sqrt(S*(S-A)*(S-B)*(S-C)).
```

- from <http://www.erlang.org/course/course.html>
- Grouping with **begin ... end**



# Modules

- Code in .hrl and .erl files
- Compilation unit: *module*
- Modules contain *forms*  
(e.g. function and macro defs)



# Example: mymath.erl

```
-module(mymath).  
  
-export([fib/1,prime/1,pi/0]).  
  
-define(PI,3.14).  
  
pi() -> ?PI.  
  
fib(N) when N<2 -> N;  
fib(N) -> fib(N-1) + fib(N-2).  
  
prime(1) -> false;  
prime(N) when N > 1 -> prime(N,2).  
  
prime(N,M) when M*M > N -> true;  
prime(N,M) when N rem M == 0 -> false;  
prime(N,M) -> prime(N,M+1).
```



# Calling functions

- Full name of exported functions:  
`mymath:prime/1`
- They can be called from other modules:  
`mymath:prime(1987)`
- They can be imported and called without module prefix
- Many built-in functions (BIFs) are auto-imported
- Calling through a variable:  
`F = fun mymath:prime/1, F(1987)`
- Dynamic call: `apply(Module, Function, Args)`



# Compiling and running

```
$ ls mymath.erl
mymath.erl
$ erl
Erlang R16B (erts-5.10.1) [source] [smp:4:4]
[async-threads:10] [hipe] [kernel-poll:false]
```

Eshell V5.10.1 (abort with ^G)

```
1> c(mymath).
{ok,mymath}
2> mymath:prime(1987).
true
3> q().
ok
4> $ ls mymath*
mymath.beam mymath.erl
```

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# What else?

- Concurrency, distributed programming
  - processes, ports, nodes
  - sending and receiving messages
- Exception handling
- Standard libraries, OTP
- Programming patterns (behaviors)



# Coming back to PaRTE...

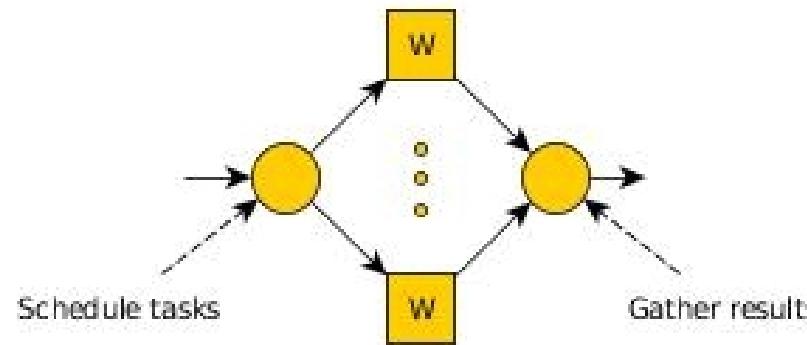
Assist parallelization of Erlang code with the

## ParaPhrase Refactoring Tool for Erlang

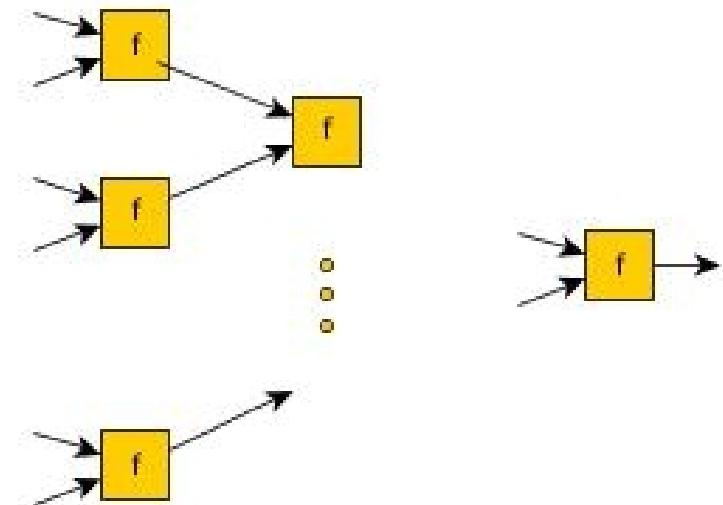
- Discover parallelizable code fragments
- Predict speedup, make suggestions
- Perform guided automated refactorings
- Pattern-based parallelism

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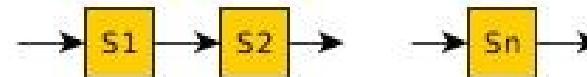
## Farm



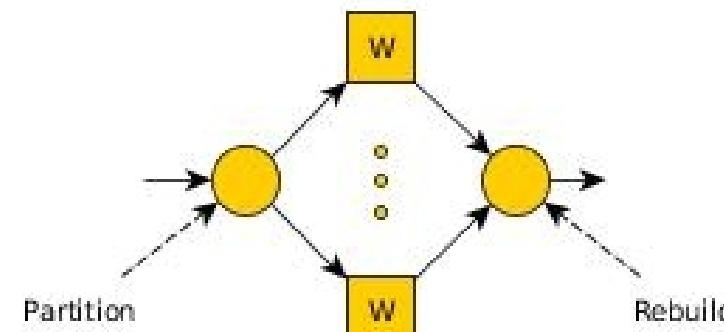
## Reduce



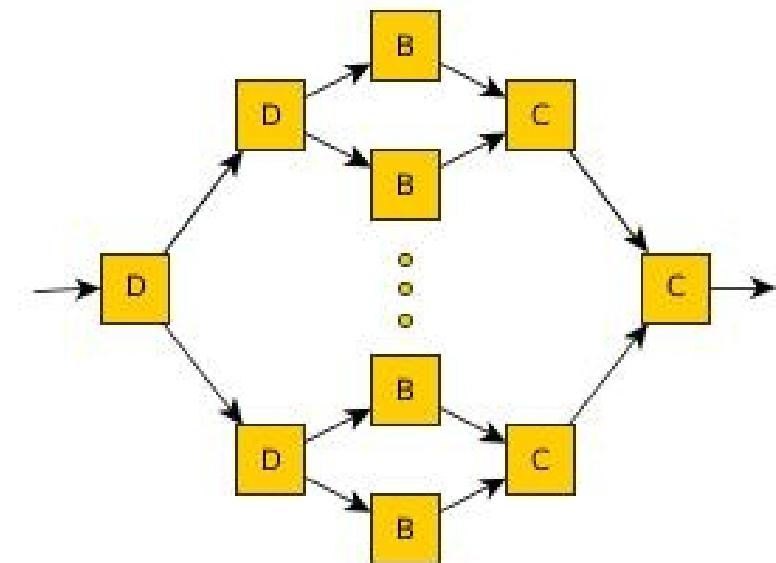
## Pipeline



## Map



## Divide&Conquer





# Parallel skeletons

<http://paraphrase-ict.eu/Deliverables/deliverable-2.6>

## The **skel** library

- Basic algorithmic skeletons  
farm, pipe, map, reduce, ord, feedback etc.
- High-level patterns: skel hlp  
dc, evolutionPool etc.
- Heterogeneous skeletons: Lapedo  
OpenCL kernels for CPU and GPU

<http://paraphrase-ict.eu/Deliverables/d27prototype.tar.gz>



# Example: parsing modules

```
[ parse ( scan ( read ( Module ) ) )
    || Module <- Modules ]
```

# Example: parsing modules

```
[ parse ( scan ( read ( Module ) ) )
    || Module <- Modules ]
```

```
skel:do([
  { farm, [{ pipe, [ { seq, fun read/1 },
                  { seq, fun scan/1 },
                  { seq, fun parse/1 }
                ] }
  ], 5 }
], Modules )
```



# Example: radix sort

```
sort( List ) -> sort(List,0).
```

```
sort( List, _ ) when length(List) < 2 ->  
List;
```

```
sort( List, Level ) ->  
lists:append(  
[ sort( Bucket, Level+1 )  
|| Bucket <- divide( List, Level )  
]  
).
```

```
divide( List, Level ) -> ...
```



# Divide-and-conquer pattern

```
{ dc, IsBase, BaseFun, Divide, Combine }
```

```
{ dc, IsBase, BaseFun, Divide, Combine,  
MaxProcesses }
```

```
SEQ_DC =  
{ seq_dc, IsBase, BaseFun, Divide, Combine },
```

```
PAR_DC =  
{ dc, IsSeq, SEQ_DC, Divide, Combine }
```



# Example: radix sort

```
sort( List ) -> skel:do( [{} dc,  
    fun({Lst,Level}) -> length(Lst) < 2 end,  
    fun({Lst,Level}) -> Lst end,  
    fun({Lst,Level}) ->  
        [ {Bucket,Level+1}  
          || Bucket <- divide(Lst, Level)  
        ]  
    end,  
    fun lists:append/1  
    [] , {List,0} ).
```





# Example: radix sort

```
sort( List ) -> sk_hlp:dc(  
    fun({Lst,Level}) -> length(Lst) < 2 end,  
    fun({Lst,Level}) -> Lst end,  
    fun({Lst,Level}) ->  
        [ {Bucket,Level+1}  
          || Bucket <- divide(Lst, Level)  
        ]  
    end,  
    fun lists:append/1  
        ) ({List,0}).
```



# ParaPhrase approach

- Identify (strongly hygienic) components
- Discover patterns of parallelism
- Structure the components into a parallel program
  - Turn the patterns into concrete code (skeletons)
  - Take performance, energy etc. into account
- Restructure if necessary
- Use a refactoring tool

# Effectiveness of the approach

Parallelization	Manual	ParaPhrasing
Convolution	3 days	3 hours
Ant Colony	1 day	1 hour
BasicN2	5 days	5 hours
Graphical Lasso	12 hours	2 hours

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# PaRTE

## ParaPhrase Refactoring Tool for Erlang

- Locate parallel pattern candidates
- Estimate speedup for different configurations
- Advise programmer
- Assist with refactoring
- Enforce preservation of functionality

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# Googling the code for patterns

## Pattern Candidate Browser

### Transformation sequences

ID	Configuration	Module	Function	Arity	Number of workers	Expected speedup (CPU)	Expected speedup (GPU)	Recommended?
1 ( $\Delta e295$ )		matrix_ex_paper	mult_matrix2	2	12	11,99	1,00	✓
2 ( $\Delta e243$ )		matrix_ex_paper	mult_matrix	2	12	10,80	1,00	✓
6 ( $\Delta(\Delta e337)$ )		matrix_ex_paper	mult_matrix2	2	12	6,58	1,00	✓
3 ( $\Delta(\Delta e337)$ )		matrix_ex_paper	mult_matrix	2	12	6,58	1,00	✓
5 ( $\Delta e292$ )		matrix_ex_paper	mult_matrix2	2	12	2,98	1,00	✓
4 ( $\Delta e337$ )		matrix_ex_paper	scalar_product	2	12	1,06	1,00	✓

Chart options ▾

Apply selected transformations

### Details of the transformation sequence

Configuration	Location information	Program text	Number of workers	Sequential CPU time	Sequential GPU time	Parallel CPU time	Parallel GPU time	Expected speedup (CPU)	Expected speedup (GPU)	Used stream length
e337	/Users/V/paraphrase/referl/tool/matrix/matrix_ex_paper.erl : {{18,15},{18,25}} - {{18, 30}, {18, 30}}	mult_scalar(A,B)	1	0,14	0,00	0,14	0,00	1,00	1,00	1
( $\Delta e337$ )	/Users/V/paraphrase/referl/tool/matrix/matrix_ex_paper.erl [ mult_scalar(A,B)    : {{18,13},{18,13}} - {{19, 39}, {19, 39}} {A,B} <- lists:zip(R,C) ]		1	1 375,42	0,00	2 506,26	0,00	0,55	1,00	10 000
( $\Delta(\Delta e337)$ )	/Users/V/paraphrase/referl/tool/matrix/matrix_ex_paper.erl [ scalar_product(R,C)    R : {{6,3},{6,3}} - {{7, 26}, {7, 26}} <- Rows, C <- Cols ]		12	13 754 154,08	0,00	2 091 407,67	0,00	6,58	1,00	10 000

Chart options ▾

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# Candidate Browser

## Transformation sequences

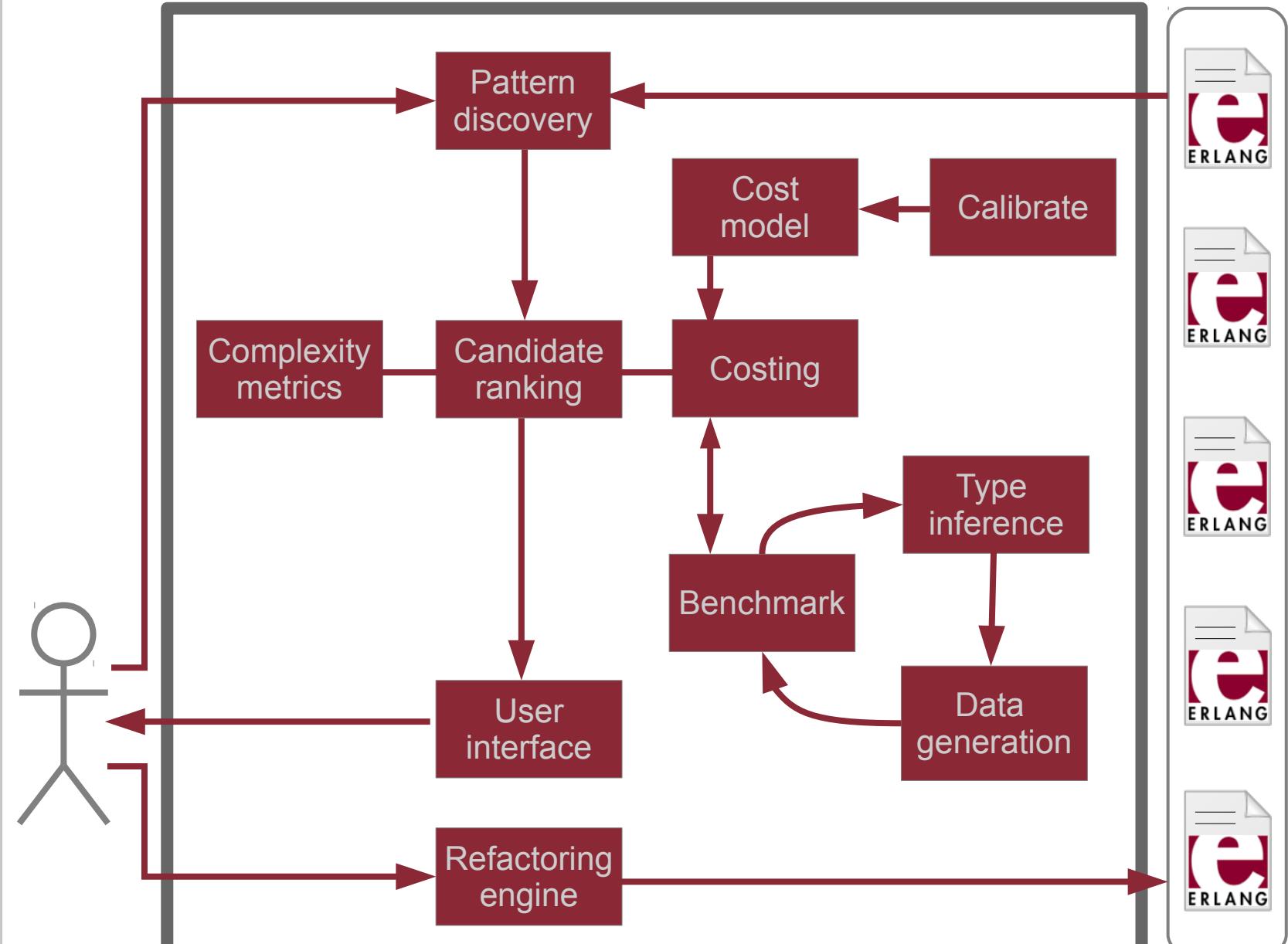
Configuration	Module	Function	Arity	Number of workers	Expected speedup (CPU)
(Δe295)	matrix_ex_paper	mult_matrix2	2	12	11,99
(Δe243)	matrix_ex_paper	mult_matrix	2	12	10,80
(Δ(Δe337))	matrix_ex_paper	mult_matrix2	2	12	6,58
(Δ(Δe337))	matrix_ex_paper	mult_matrix	2	12	6,58
(Δe292)	matrix_ex_paper	mult_matrix2	2	12	2,98
(Δe337)	matrix_ex_paper	scalar_product	2	12	1,06

Sort options ▾

## Details of the transformation sequence

Configuration	Location information	Program text	Number of workers	Sequential CPU time	Sequential GPU time	Parallel CPU time
	/Users/V/paraphrase/referl/tool/matrix/matrix_ex_paper.erl : {{18,15},{18,25}} - {{18, 30}, {18, 30}}	mult_scalar(A,B)	1	0,14	0,00	
7)	/Users/V/paraphrase/referl/tool/matrix/matrix_ex_paper.erl : {{18,13},{18,13}} - {{19, 39}, {19, 39}}	[ mult_scalar(A,B)    {A,B} <- lists:zip(R,C) ]	1	1 375,42	0,00	2
(Δe337))	/Users/V/paraphrase/referl/tool/matrix/matrix_ex_paper.erl : {{6,3},{6,3}} - {{7, 26}, {7, 26}}	[ scalar_product(R,C)    R <- Rows, C <- Cols ]	12	13 754 154,08	0,00	2 091

Sort options ▾





# Pattern candidate discovery

- Collect syntactic & semantic information
  - List comprehensions
  - Library calls (`lists:map/2`)
  - Recursion structure
  - Side conditions
- Task farm, pipeline, divide-and-conquer, feedback
- Heuristics



# Task farm

```
[ parse(scan(read( Module )))  
  || Module <- Modules ]
```

```
Work = fun(Module) ->  
        parse(scan(read(Module))) end,  
skel:do([{farm,[{seq,Work}],5}], Modules)
```



# Pipeline

```
[ parse(scan(read( Module )))  
  || Module <- Modules ]
```

```
Stages =  
  [{seq, read/1}, {seq, scan/1}, {seq, parse/1}],  
 skel:do([{pipe, Stages}], Modules)
```



# Farm of pipes

```
[ parse(scan(read( Module )))  
  || Module <- Modules ]
```

```
Stages =  
  [{seq,read/1},{seq,scan/1},{seq,parse/1}],  
Work = {pipe,Stages},  
skel:do([{farm,Work,5}], Modules)
```



# Functional “quicksort”: d&c

```
qs ( List ) ->
  case List of
    []      -> [] ;
    [H|T]  ->
      {List1,List2} = lists:partition(
                      fun(X) -> X < H end,
                      T),
      qs(List1) ++ [H] ++ qs(List2)
  end.
```



# Karatsuba: d&c

**karatsuba( Num1 , Num2 ) ->**

...

**Z0 = karatsuba( Low1 , Low2 ),**

**Z1 = karatsuba( add(Low1,High1),**  
**add(Low2,High2) ),**

**Z2 = karatsuba( High1 , High2 ),**

...



# Radix sort: d&c

```
sort( [] , _ ) -> [] ;  
sort( [V] , _ ) -> [V] ;  
sort( List, Level ) ->  
    Buckets = divide( List, Level ),  
    SortedLists =  
        lists:map( fun(B) -> sort(B,Level+1) end,  
                  Buckets ),  
        lists:append( SortedLists ).
```

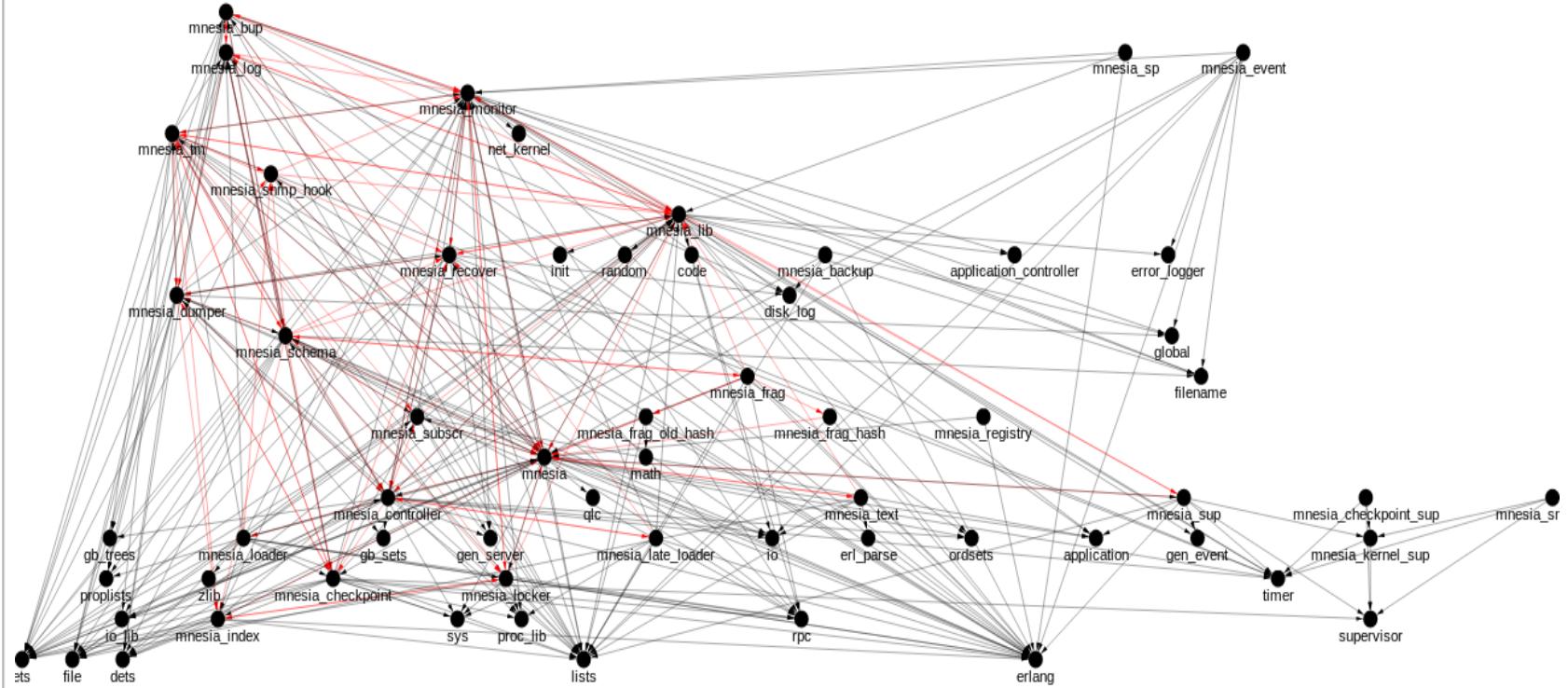
# Minimax: d&c

```
mm_max( Node, Depth ) ->  
  case Depth == 0 orelse terminal(Node)  
    true  -> value ( Node ) ;  
    false -> lists:max( [ mm_min(C, Depth-1)  
                           || C <- children(Node)  
                         ] )  
  
end .
```

```
mm_min( Node, Depth ) -> ... mm_max ...
```

# RefactorErl

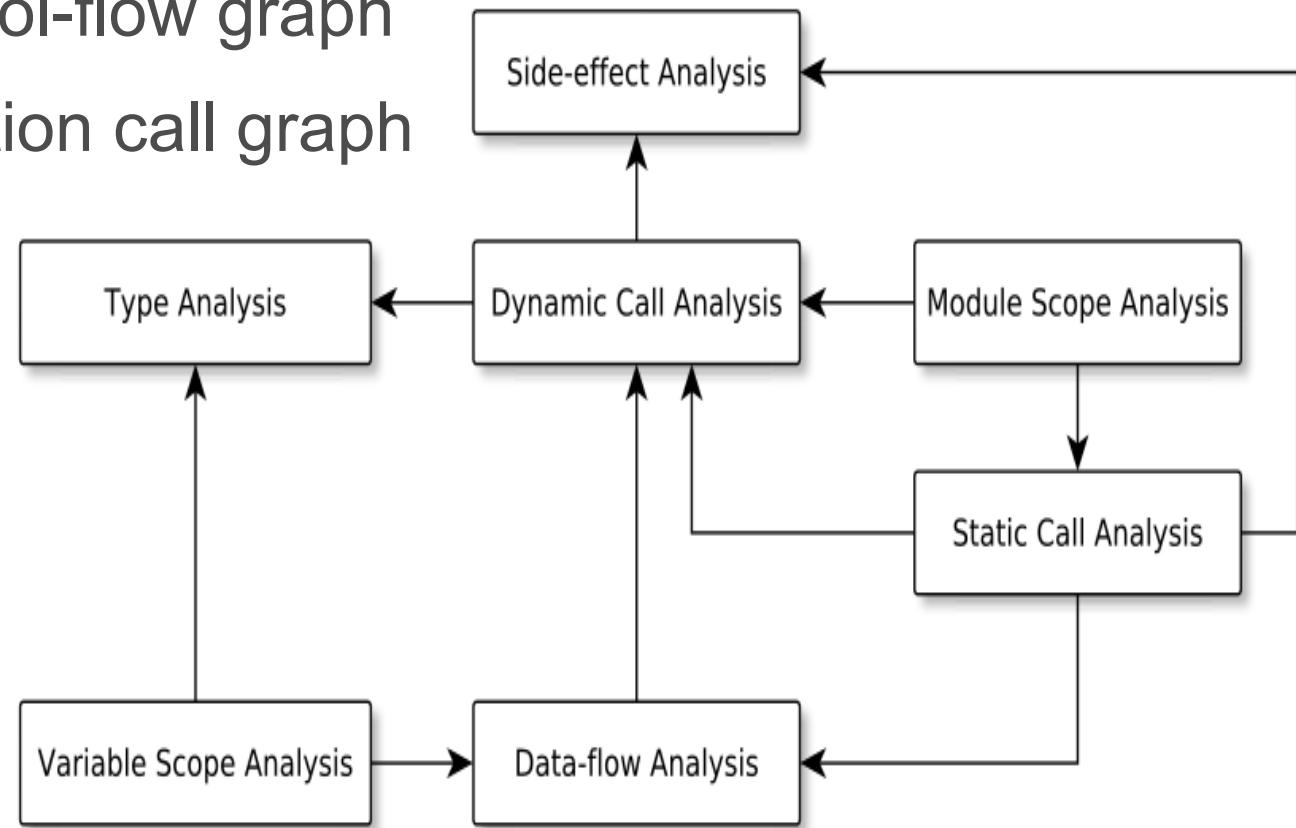
Static source code analyzer and transformer  
<http://plc.inf.elte.hu/erlang>



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# Standard static analyses

- Data-flow graph
- Control-flow graph
- Function call graph





# Pattern-specific analyses

- Identify components
  - Side effects
- Identify patterns
  - Data dependency



# Component

Action performed by Worker (farm) or Stage (pipe)

- Side-effect analysis
  - Message passing
  - NIFs and global variables
  - ETS etc.
  - Process dictionary, node names
  - Exceptions
- Hygiene rather than purity



# Hygienic component

- Identify used resources
- Classify read/alter operations

$\text{use}(C, R) \in \{ \text{No}, \text{Read}, \text{Alter} \}$

- Component set:  
components executed in parallel

$\forall R \quad \forall C_1 \neq C_2 \in S:$

$\text{use}(C_1, R) = \text{Alter} \rightarrow \text{use}(C_2, R) = \text{No}$



# Element-wise processing

```
[ parse(scan(read( Module )))  
  || Module <- Modules ]
```

# Element-wise processing

---

```
[ parse(scan(read( Module )))  
  || Module <- Modules ]
```

---

```
Work = fun(Module) ->  
        parse(scan(read(Module))) end,  
lists:map(Work,Modules)
```

---

```
psr([]) -> [];  
psr([H|T]) -> [parse(scan(read(H))) | psr(T)].  
... psr(Modules) ...
```

---



# Element-wise processing

```
f(P1, P2, P3, P4) ->
  case P3 of
    [] -> [];
    [ Head | Tail ] ->
      X = ... Head ... ,
      [ X | f(P1,P2,Tail,P4) ]
  end.
```

a map-like function

# Element-wise processing

```
f(P1, P2, P3, P4) ->
  case P3 of
    [] -> [];
    [ Head | Tail ] ->
      X = ... Head ... ,
      [ X | f(P1,P2,Tail,P4) ]
  end.
```

a map-like function

- $f$  must be recursive:  
the interprocedural CFG must contain an execution path  
from the “starting node” of  $f$  to a “call-node” of  $f$   
 $\exists p \in EP(start_f) \text{ such that } call_f \in p$
- ... base case ... single recursive call ... regularities ...  
... data dependencies ... compact data flow reaching ...



# Divide-and-conquer

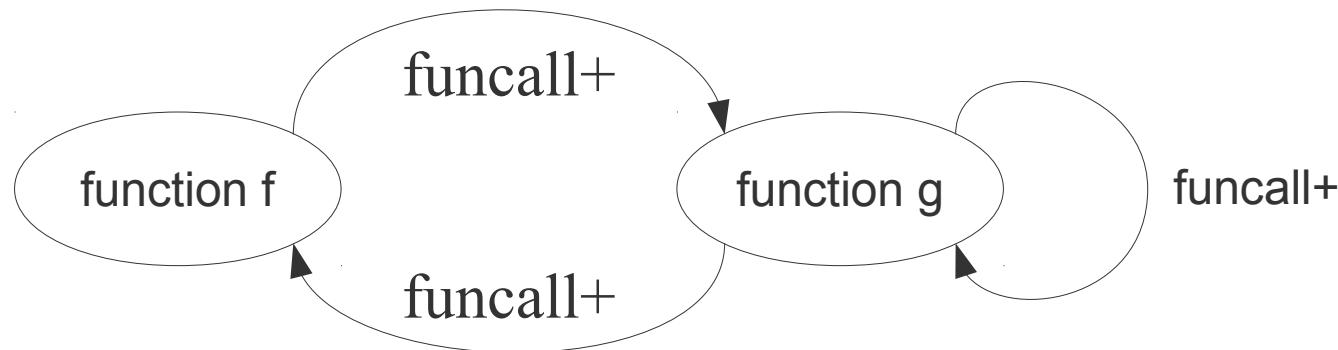
- A function has multiple recursive calls to itself
- The arguments of a recursive call do not depend on the result of another recursive call

Costly!

- *f must be recursive:*
  - the interprocedural CFG must contain an execution path from the “starting node” of *f* to a “call-node” of *f* $\exists p \in EP(start_f) \text{ such that } call_f \in p$
  - etc.

# Faster rules

- If  $h$  operates on lists, and
  - contains list comprehension with  $h$  in head
  - passes  $h$  to *lists:map/2* or a map-like function
- Analyze the function call graph





# Experiments

- Ant Colony Optimization  
Single Machine Total Weighted Tardiness Problem
- Image merging case study
- Intensional Computing Engine  
Evaluator of the abstract syntax tree of ICE
- Evolutionary Multi-Agent Systems framework
- Thorn  
Map-reduce framework
- Mnesia  
Distributed database management system
- RefactorErl core  
Static program analysis&transformation framework

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# Problem sizes

	ELOC	Functions	Files
Ant Colony Optimization	483	56	21
Image merging	779	104	6
ICE evaluator	1094	141	21
Thorn	1313	158	15
EMAS framework	1646	177	25
RefactorErl core	19694	1534	53
Mnesia	22653	1693	31



# Discovery results

	farm	pipe	reduce	d&c	pool
Ant Colony Optimization	10				
Image merging	34	4	2		
ICE evaluator	7		4		
Thorn	17		5		
EMAS framework	71	20	16		9
RefactorErl core	486	49	55	31	
Mnesia	135	8	36	57	2

Map-like functions: 9

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# Some interesting candidates

- map-like functions
- divide-and-conquer algorithms

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# Map-like function (Mnesia)

```
reverse([]) -> [];  
  
reverse([ H=#commit{ ram_copies      = Ram,  
                    disc_copies     = DC,  
                    disc_only_copies= DOC,  
                    snmp           = Snmp }  
        | R ]) ->  
    [ H#commit{  
        ram_copies      = lists:reverse(Ram),  
        disc_copies     = lists:reverse(DC),  
        disc_only_copies= lists:reverse(DOC),  
        snmp           = lists:reverse(Snmp)  
    }  
    | reverse(R) ].
```



# Map-like function (EMAS)

```
count_funstats(_, []) -> [];  
count_funstats(  
    Agents,  
    [{Stat, MapFun, ReduceFun, OldAcc} | T]  
) ->  
    NewAcc =  
        lists:foldl( ReduceFun, OldAcc,  
            [MapFun(Agent) || Agent <- Agents] ),  
        [ {Stat, MapFun, ReduceFun, NewAcc}  
        | count_funstats(Agents, T)  
        ].
```



# D&C (RefactorErl)

...

```
listcons_length(N, #expr{}) ->  
    Ns = ?Dataflow:?reach([N], [back], true),  
    L1 = [N2 || N2 <- Ns, N2 /= N,  
          ?Graph:class(N2) == expr],  
    {L2, L3} = lists:partition(  
        fun is_cons_expr/1, L1 ),  
    if L2 == [] orelse L3 /= [] ->  
        incalculable;  
    true ->  
        lists:append(lists:map(  
            fun listcons_length/1, L2 ))  
    end;
```

...

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# Another D&C (RefactorErl)

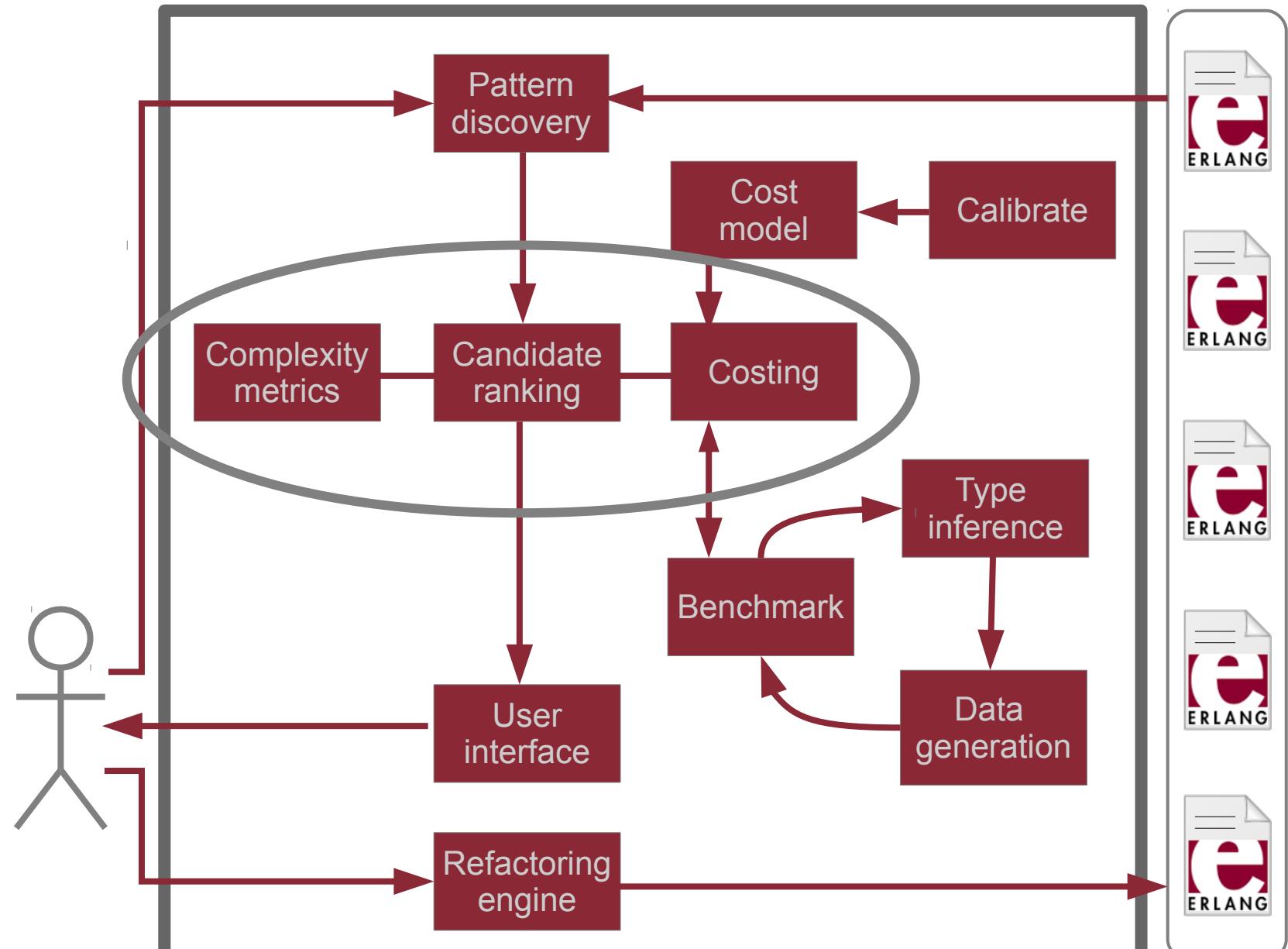
```

realtoken_neighbour(Node, DirFun, DownFun) ->
    case lists:member(?Graph:class(Node),[clause,expr,form,typexp,lex]) of
        false -> no;
        _ -> case ?Syn:parent(Node) of
            [] -> no;
            [{_,Parent}] -> case lists:dropwhile( fun({_T,N}) -> N/=Node end,
                                                DirFun(?Syn:children(Parent)))
                                of
                                    [{_,Node},{_,NextNode}|_] -> DownFun(NextNode);
                                    _ -> realtoken_neighbour(Parent, DirFun, DownFun)
            end;
            Parents -> realtoken_neighbour_( Parents, DownFun(Node),
                                         DirFun, DownFun )
    end
end.

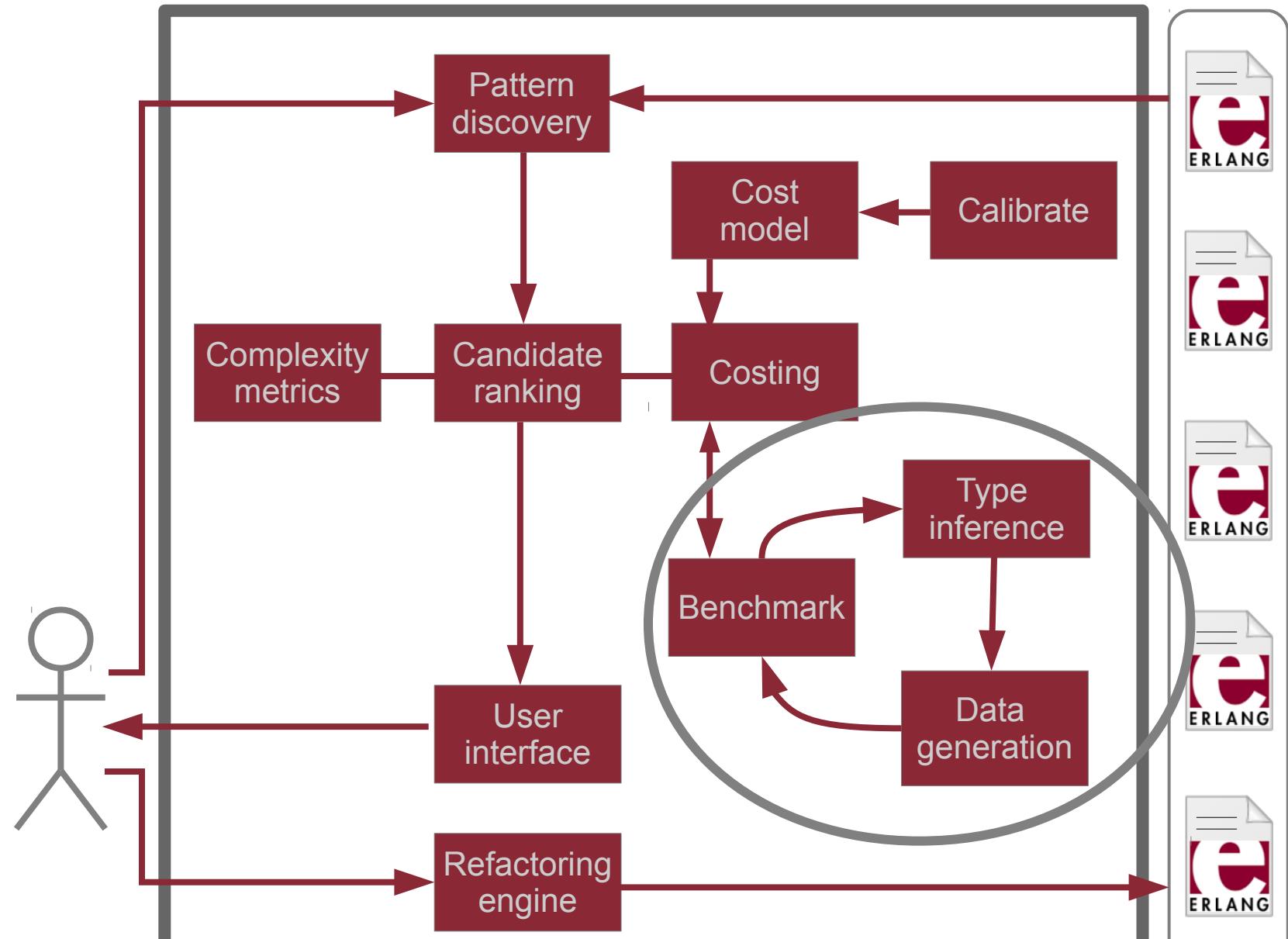
% Implementation helper function for realtoken_neighbour/3

realtoken_neighbour_([], _FirstLeaf,_DirFun,_DownFun) -> no;
realtoken_neighbour_([{_,Parent}|Parents], FirstLeaf, DirFun, DownFun) ->
    case realtoken_neighbour(Parent, DirFun, DownFun) of
        FirstLeaf -> realtoken_neighbour_(Parents, FirstLeaf, DirFun, DownFun);
        NextLeaf -> NextLeaf
    end.

```



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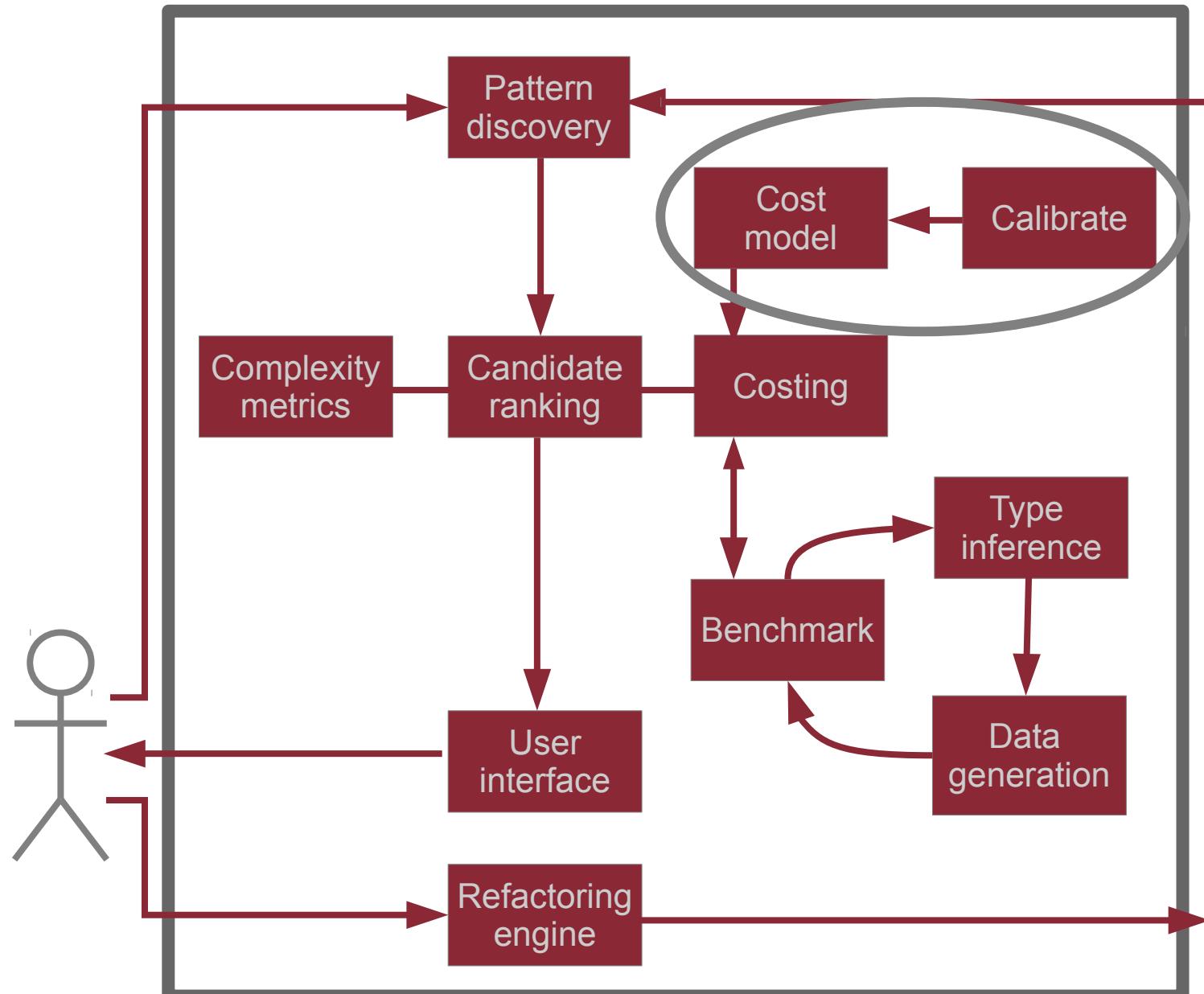
# Benchmarking

- Split up pattern candidates into components
- Determine free variables (inputs)
- Assemble a new module
  - Components turned into functions
  - Instrumented with time measurements
- Load module
- Generate random input and profile
- Make statistics



# Random input?

- Not always meaningful...
- ... but easy to automate!
- Find out the type of free variables (TypEr)
- QuickCheck generates values by type



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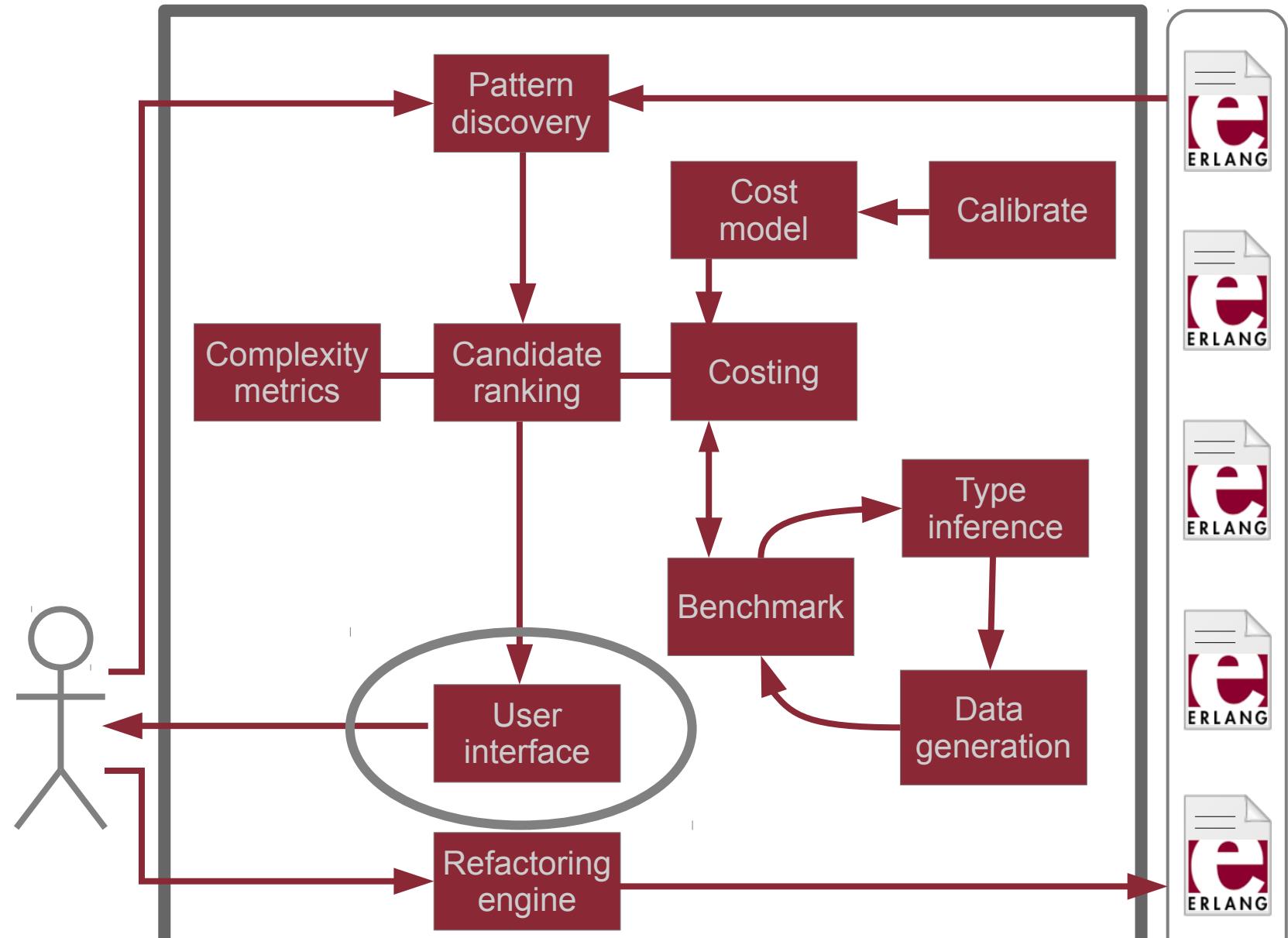


# Cost model

- Approximation
- E.g. for farm:

$$\begin{aligned} T_{farm} := & T_{work} * \lceil L / \min(N_p, N_w) \rceil + \\ & T_{spawn} * (N_w + 2) + \\ & T_{copy}(L) * 3 + T_{spawn} + T_{copy}(L) * 2 \end{aligned}$$

- 
- Needs calibration!



# Pattern Candidate Browser

- After ranking pattern candidates
- Web-based interface
  - Information for decision making
  - Not too many details
- Services
  - Multiple users
  - Persistent results
  - Export XML, JSON, CSV, Erlang terms

Pattern Candidate Browser

Transformation sequences

ID	Configuration	Module	Function	Arity	Number of workers	Expected speedup (CPU)	Expected speedup (GPU)	Recommended?
1 (de295)		matrix_ex_paper	mult_matrix2	2	12	11,99	1,00	✓
2 (de243)		matrix_ex_paper	mult_matrix	2	12	10,80	1,00	✓
6 (Δde337))		matrix_ex_paper	mult_matrix2	2	12	6,58	1,00	✓
3 (Δde337))		matrix_ex_paper	mult_matrix	2	12	6,58	1,00	✓
5 (de292)		matrix_ex_paper	mult_matrix2	2	12	2,98	1,00	✓
4 (de337)		matrix_ex_paper	scalar_product	2	12	1,06	1,00	✓

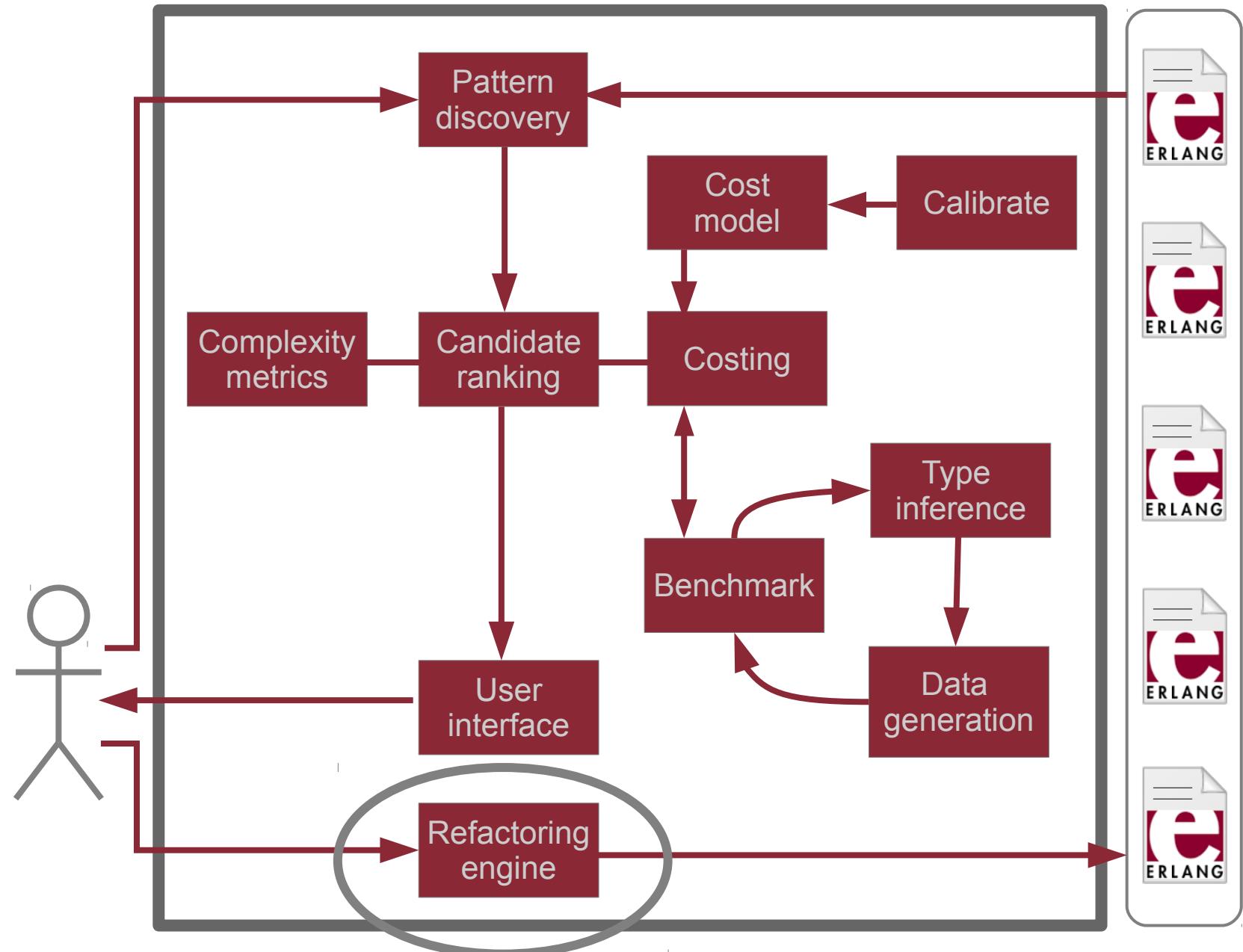
Chart options      Apply selected transformations

Details of the transformation sequence

Configuration	Location information	Program text	Number of workers	Sequential CPU time	Sequential GPU time	Parallel CPU time	Parallel GPU time	Expected speedup (CPU)	Expected speedup (GPU)	Used stream length
e337	/Users/V/paraphrase/refinertool/matrix/matrix_ex_paper.erl	mult_scaler(A,B) : {t(18,15),(18,25)} -> {t(18,30),(18,30)}	1	0,14	0,00	0,14	0,00	1,00	1,00	1
(Δde337)	/Users/V/paraphrase/refinertool/matrix/matrix_ex_paper.erl	[ mult_scaler(A,B)    {t(18,13),(18,13)} -> {t(18,39),(19,39)} ]	1	1 375,42	0,00	2 506,26	0,00	0,55	1,00	10 000
(Δ(de337))	/Users/V/paraphrase/refinertool/matrix/matrix_ex_paper.erl	[ scalar_product(R,C)    R < Rows, C < Cols ] : {t(8,3),(8,3)} -> {t(7,26),(7,26)}	12 13 754	154,08	0,00	2 091 407,67	0,00	6,58	1,00	10 000

Chart options

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# Refactoring

- Program Shaping
- Introduction of Skeletons
- Cleanup Transformations

Either directly or after PC discovery



# Program shaping

- Restructures and tunes program to get it into appropriate shape for skeleton introduction
- Removal of:
  - Dependencies
  - Locks
  - Global State
  - Nestings
  - Copying



# Future work: more refactorings

```
sort( List ) -> sort(List,0).
sort( List, _ ) when length(List) < 2 -> List;
sort( List, Level ) ->
    lists:append(
        [sort( Bucket, Level+1 ) || Bucket <- divide( List, Level )]
    ).
```



```
sort( List ) -> skel:do( [{ dc,
    fun({Lst,Level}) -> length(Lst) < 2 end,
    fun({Lst,Level}) -> Lst end,
    fun({Lst,Level}) ->
        [ {Bucket,Level+1} || Bucket <- divide(Lst, Level) ]
    end,
    fun lists:append/1
}], {List,0}).
```

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# Resources

ParaPhrase project (FP7 contract no. 288570)  
<http://paraphrase-ict.eu/>

ParaPhrase @ ELTE  
<http://paraphrase-enlarged.elte.hu/>

PaRTE docs & download:  
<http://pnyf.inf.elte.hu/trac/refactorerl/wiki/partे>

CEFP 2015 instructions:  
<http://pnyf.inf.elte.hu/trac/refactorerl/wiki/partе/cefp>

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# Conclusions

- Pattern-based parallelism
- **ParaPhrase Refactoring Tool for Erlang**

Pattern discovery and refactoring

Candidates prioritized,  
support for decision making

- Finds many places to introduce parallelism
- Supports candidate ranking
- Works effectively with a smart programmer
- Offers performance gains with small effort